

COMP4418: Knowledge Representation and Reasoning—Solutions to Exercise 1

Propositional Logic

1. (i) $(\neg Ja \wedge \neg Jo) \rightarrow T$

Where:

Ja: Jane is in town

Jo: John is in town

T: we will play tennis

(ii) $R \vee \neg R$

Where:

R: it will rain today

(iii) $\neg S \rightarrow \neg P$

Where:

S: you study

P: you will pass this course

2. (i) $P \rightarrow Q$

$\neg P \vee Q$ (remove \rightarrow)

(ii) $(P \rightarrow \neg Q) \rightarrow R$

$\neg(\neg P \vee \neg Q) \vee R$ (remove \rightarrow)

$(\neg\neg P \wedge \neg\neg Q) \vee R$ (De Morgan)

$(P \wedge Q) \vee R$ (Double Negation)

$(P \vee R) \wedge (Q \vee R)$ (Distribute \vee over \wedge)

(iii) $\neg(P \wedge \neg Q) \rightarrow (\neg R \vee \neg Q)$

$\neg\neg(P \wedge \neg Q) \vee (\neg R \vee \neg Q)$ (remove \rightarrow)

$(P \wedge \neg Q) \vee (\neg R \vee \neg Q)$ (Double Negation)

$(P \vee \neg R \vee \neg Q) \wedge (\neg Q \vee \neg R \vee \neg Q)$ (Distribute \vee over \wedge)

This can be further simplified to: $((P \vee \neg R \vee \neg Q) \wedge (\neg Q \vee \neg R))$

And in fact this can be simplified to $\neg Q \vee \neg R$ since $(\neg Q \vee \neg R) \vdash$

$(P \vee \neg R \vee \neg Q)$

3. (i)

| <i>P</i> | <i>Q</i> | $P \rightarrow Q$ | $\neg Q$ | $\neg P$ |
|----------|----------|-------------------|----------|----------|
| <i>T</i> | <i>T</i> | <i>T</i> | <i>F</i> | <i>T</i> |
| <i>T</i> | <i>F</i> | <i>F</i> | <i>T</i> | <i>F</i> |
| <i>F</i> | <i>T</i> | <i>T</i> | <i>F</i> | <i>T</i> |
| <i>F</i> | <i>F</i> | <i>T</i> | <i>T</i> | <i>T</i> |

In all rows where both $P \rightarrow Q$ and $\neg Q$ are true, $\neg P$ is also true.

Therefore, inference is valid.

(ii)

| <i>P</i> | <i>Q</i> | $\neg P$ | $\neg Q$ | $P \rightarrow Q$ | $\neg Q \rightarrow \neg P$ |
|----------|----------|----------|----------|-------------------|-----------------------------|
| <i>T</i> | <i>T</i> | <i>F</i> | <i>F</i> | <i>T</i> | <i>T</i> |
| <i>T</i> | <i>F</i> | <i>F</i> | <i>T</i> | <i>F</i> | <i>F</i> |
| <i>F</i> | <i>T</i> | <i>T</i> | <i>F</i> | <i>T</i> | <i>T</i> |
| <i>F</i> | <i>F</i> | <i>T</i> | <i>T</i> | <i>T</i> | <i>T</i> |

In all rows where both $P \rightarrow Q$ is true, $\neg Q \rightarrow \neg P$ is also true.

Therefore, inference is valid.

| P | Q | R | $P \rightarrow Q$ | $Q \rightarrow R$ | $P \rightarrow R$ |
|-----|-----|-----|-------------------|-------------------|-------------------|
| T | T | T | T | T | T |
| T | T | F | T | F | F |
| T | F | T | F | T | T |
| T | F | F | F | T | F |
| F | T | T | T | T | T |
| F | T | F | T | F | T |
| F | F | T | T | T | T |
| F | F | F | T | T | T |

In all rows where both $P \rightarrow Q$ and $Q \rightarrow R$ are true, $P \rightarrow R$ is also true. Therefore, inference is valid.

4. (i) $\text{CNF}(P \rightarrow Q)$
 $\equiv \neg P \vee Q$

$\text{CNF}(\neg Q)$
 $\equiv \neg Q$

$\text{CNF}(\neg\neg P)$
 $\equiv P$ (Double Negation)

Proof:

1. $\neg P \vee Q$ (Hypothesis)
2. $\neg Q$ (Hypothesis)
3. P (Negation of Conclusion)
4. Q 1, 3 Resolution
5. \square 2, 4 Resolution

(ii) $\text{CNF}(P \rightarrow Q)$
 $\equiv \neg P \vee Q$

$\text{CNF}(\neg(\neg Q \rightarrow \neg P))$
 $\equiv \neg(\neg\neg Q \vee \neg P)$ (Remove \rightarrow)
 $\equiv \neg(Q \vee \neg P)$ (Double Negation)
 $\equiv \neg Q \wedge \neg\neg P$ (De Morgan)
 $\equiv \neg Q \wedge P$ (Double Negation)

Proof:

1. $\neg P \vee Q$ (Hypothesis)
2. $\neg Q$ (Negation of Conclusion)
3. P (Negation of Conclusion)
4. $\neg P$ 1, 2 Resolution
5. \square 3, 4 Resolution

(iii) $P \rightarrow Q, Q \rightarrow R \vdash P \rightarrow R$

$\text{CNF}(P \rightarrow Q)$
 $\equiv \neg P \vee Q$

$\text{CNF}(Q \rightarrow R)$
 $\equiv \neg Q \vee R$

$$\begin{aligned}
& \text{CNF}(\neg(P \rightarrow R)) \\
& \equiv \neg(\neg P \vee R) \text{ (Remove } \rightarrow) \\
& \equiv \neg\neg P \wedge \neg R \text{ (De Morgan)} \\
& \equiv P \wedge \neg R \text{ (Double Negation)}
\end{aligned}$$

Proof:

1. $\neg P \vee Q$ (Hypothesis)
2. $\neg Q \vee R$ (Hypothesis)
3. P (Negation of Conclusion)
4. $\neg R$ (Negation of Conclusion)
5. Q 1, 3 Resolution
6. R 2, 5 Resolution
7. \square 4, 6 Resolution

5. (i)

| P | Q | $\neg P$ | $P \vee Q$ | $(P \vee Q) \wedge \neg P$ | $((P \vee Q) \wedge \neg P) \rightarrow Q$ |
|-----|-----|----------|------------|----------------------------|--|
| T | T | F | T | F | T |
| T | F | F | T | F | T |
| F | T | T | T | T | T |
| F | F | T | F | F | T |

Last column is always true no matter what truth assignment to the atoms P and Q . Therefore $((P \vee Q) \wedge \neg P) \rightarrow Q$ is a tautology.

- (ii) $((P \rightarrow Q) \wedge \neg(P \rightarrow R)) \rightarrow (P \rightarrow Q)$

(iii)

| P | Q | R | $P \rightarrow Q$ | $P \rightarrow R$ | $(P \rightarrow Q) \wedge \neg(P \rightarrow R)$ | $P \rightarrow Q$ | $((P \rightarrow Q) \wedge \neg(P \rightarrow R)) \rightarrow (P \rightarrow Q)$ |
|-----|-----|-----|-------------------|-------------------|--|-------------------|--|
| T | T | T | T | T | T | T | T |
| T | T | F | T | F | F | F | T |
| T | F | T | F | T | F | T | T |
| T | F | F | F | T | F | T | T |
| F | T | T | T | T | T | T | T |
| F | T | F | T | F | F | F | T |
| F | F | T | T | T | T | T | T |
| F | F | F | T | T | T | T | T |

Last column is always true no matter what truth assignment to the atoms P , Q and R . Therefore $((P \rightarrow Q) \wedge \neg(P \rightarrow R)) \rightarrow (P \rightarrow Q)$ is a tautology.

(iv)

| P | $\neg P$ | $\neg P \wedge P$ | $\neg(\neg P \wedge P)$ | $\neg(\neg P \wedge P) \wedge P$ |
|-----|----------|-------------------|-------------------------|----------------------------------|
| T | F | F | T | T |
| F | T | F | T | F |

Last column is not always true. Therefore $\neg(\neg P \wedge P) \wedge P$ is not a tautology.

- (v) $(P \vee Q) \rightarrow \neg(\neg P \wedge \neg Q)$

| P | Q | $\neg P$ | $\neg Q$ | $P \vee Q$ | $\neg P \wedge \neg Q$ | $\neg(\neg P \wedge \neg Q)$ | $(P \vee Q) \rightarrow \neg(\neg P \wedge \neg Q)$ |
|-----|-----|----------|----------|------------|------------------------|------------------------------|---|
| T | T | F | F | T | F | T | T |
| T | F | F | T | T | F | T | T |
| F | T | T | F | T | F | T | T |
| F | F | T | T | F | T | F | T |

6. (i) $\text{CNF}(\neg((P \vee Q) \wedge \neg P) \rightarrow Q) \equiv \neg(\neg((P \vee Q) \wedge \neg P) \vee Q)$ (Remove \rightarrow)
 $\equiv \neg\neg((P \vee Q) \wedge \neg P) \wedge \neg Q$ (DeMorgan)
 $\equiv (P \vee Q) \wedge \neg P \wedge \neg Q$ (Double Negation)

Proof:

1. $P \vee Q$ (Negated Conclusion)
2. $\neg P$ (Negated Conclusion)
3. $\neg Q$ (Negated Conclusion)
4. Q 1, 2 Resolution
5. \square 3, 4 Resolution

Therefore $\neg((P \vee Q) \wedge \neg P) \rightarrow Q$ is a tautology.

- (ii) $\text{CNF}(\neg(((P \rightarrow Q) \wedge \neg(P \rightarrow R)) \rightarrow (P \rightarrow Q)))$
 $\equiv \neg(\neg((\neg P \vee Q) \wedge \neg(\neg P \vee R)) \vee (\neg P \vee Q))$ (Remove \rightarrow)
 $\equiv \neg\neg((\neg P \vee Q) \wedge \neg(\neg P \vee R)) \wedge \neg(\neg P \vee Q)$ (De Morgan)
 $\equiv (\neg P \vee Q) \wedge (\neg\neg P \wedge \neg R) \wedge (\neg\neg P \wedge \neg Q)$ (Double Negation and De Morgan)
 $\equiv (\neg P \vee Q) \wedge (P \wedge \neg R) \wedge (P \wedge \neg Q)$ (Double Negation)

Proof:

1. $\neg P \vee Q$ (Negated Conclusion)
2. P (Negated Conclusion)
3. $\neg R$ (Negated Conclusion)
4. $\neg Q$ (Negated Conclusion)
5. Q 1, 2 Resolution
6. \square 4, 5 Resolution

Therefore $((P \rightarrow Q) \wedge \neg(P \rightarrow R)) \rightarrow (P \rightarrow Q)$ is a tautology.

- (iii) $\text{CNF}(\neg(\neg(\neg P \wedge P) \wedge P))$
 $\equiv \neg\neg(\neg P \wedge P) \vee \neg P$ (De Morgan)
 $\equiv (\neg P \wedge P) \vee \neg P$ (Double Negation)
 $\equiv (\neg P \vee \neg P) \vee (P \vee \neg P)$ (Distribute \wedge over \vee)
 $\equiv \neg P$ (Can simplify to this by removing repetition and tautologies)

Proof:

1. $\neg P$ (Negated Conclusion)

Cannot obtain empty clause using resolution so $\neg(\neg P \wedge P) \wedge P$ is not a tautology.

- (iv) $\text{CNF}(\neg((P \vee Q) \rightarrow \neg(\neg P \wedge \neg Q))) \equiv \neg(\neg(P \vee Q) \vee \neg(\neg P \wedge \neg Q))$
(Remove \rightarrow)
 $\equiv \neg\neg(P \vee Q) \vee \neg\neg(\neg P \wedge \neg Q)$ (De Morgan)
 $\equiv (P \vee Q) \vee (\neg P \wedge \neg Q)$ (Double Negation)

Proof:

1. $(P \vee Q)$ (Negated Conclusion)
2. $\neg Q$ (Negated Conclusion)
3. $\neg P$ (Negated Conclusion)
4. Q 1, 2 Resolution
5. \square 3, 4, Resolution

Therefore $(P \vee Q) \rightarrow \neg(\neg P \wedge \neg Q)$ is a tautology.