10. Iterative Compression

COMP6741: Parameterized and Exact Computation

Serge Gaspers

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Contents

1 Introduction

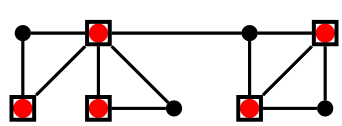
Iterative Compression
For a minimization problem:

- **Compression step:** Given a solution of size $k+1$, compress it to a solution of size k or prove that there is no solution of size k
- **Iteration step:** Incrementally build a solution to the given instance by deriving solutions for larger and larger subinstances

- Seen a lot of success in Parameterized Complexity
- Examples: best known fixed parameter algorithms for (DIRECTED) FEEDBACK VERTEX SET, EDGE BIPARTIZATION, ALMOST 2-SAT, ...

Example: Vertex Cover
A *vertex cover* in a graph $G = (V, E)$ is a subset of its vertices $S \subseteq V$ such that every edge of G has at least one endpoint in S.

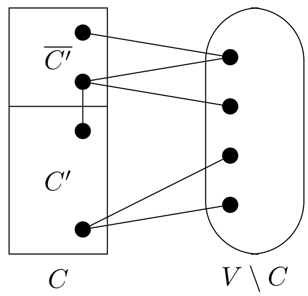
VERTEX COVER	
Input:	A graph $G = (V, E)$ and an integer k
Parameter:	k
Question:	Does G have a vertex cover of size k?



We will design a (slow) iterative compression algorithm for VERTEX COVER to illustrate the technique.

Vertex Cover: Compression Step

Comp-VC	
Input:	graph $G = (V, E)$, integer k, vertex cover C of size $k+1$ of G
Output:	a vertex cover C^* of size $\leq k$ of G if one exists



- Go over all partitions $(C', \overline{C'})$ of C
- $C^* = C' \cup N(\overline{C'})$
- If $\overline{C'}$ is an independent set and $|C^*| \leq k$ then return C^*

Vertex Cover: Iteration Step
Use algorithm for Comp-VC to solve Vertex Cover.

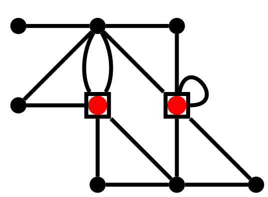
- Order vertices: $V = \{v_1, v_2, \ldots, v_n\}$
- Define $G_i = G[\{v_1, v_2, \ldots, v_i\}]$
- $C_0 = \emptyset$
- For $i = 1..n$, find a vertex cover C_i of size $\leq k$ of G_i using the algorithm for Comp-VC with input G_i and $C_{i-1} \cup \{v_i\}$. If G_i has no vertex cover of size $\leq k$, then G has no vertex cover of size $\leq k$.

Final running time: $O^*(2^k)$

2 Feedback Vertex Set

A *feedback vertex set* of a multigraph $G = (V, E)$ is a set of vertices $S \subseteq V$ such that $G - S$ is acyclic.

Feedback Vertex Set (FVS)	
Input:	Multigraph $G = (V, E)$, integer k
Parameter:	k
Question:	Does G have a feedback vertex set of size at most k?



Note: We already saw an $O^*((3k)^k)$ time algorithm for FVS. We will now aim for a $O^*(c^k)$ time algorithm, with $c \in O(1)$.

Compression Problem

Comp-FVS	
Input:	graph $G = (V, E)$, integer k, feedback vertex set S of size $k+1$ of G
Output:	a feedback vertex set S^* of size $\leq k$ of G if one exists

Iteration step

- Order vertices: $V = \{v_1, v_2, \ldots, v_n\}$
- Define $G_i = G[\{v_1, v_2, \ldots, v_i\}]$
- $S_0 = \emptyset$
- For $i = 1..n$, find a feedback vertex set S_i of size $\leq k$ of G_i using the algorithm for Comp-FVS with input G_i and $S_{i-1} \cup \{v_i\}$. If G_i has no feedback vertex set of size $\leq k$, then G has no feedback vertex set of size $\leq k$.

Suppose Comp-FVS can be solved in $O^*(c^k)$ time. Then, using this iteration, FVS can be solved in $O^*(c^k)$ time.

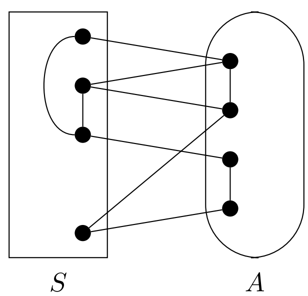
Compression step
To solve Comp-FVS, go through all partitions $(S', \overline{S'})$ of S. For each of them, we will want to find a feedback vertex set S^* of G with $|S^*| < |S|$ and $S' \subseteq S^* \subseteq V \setminus \overline{S'}$ if one exists. Equivalently, find a feedback vertex set S'' of $G - S'$ with $|S''| < |\overline{S'}|$ and $S'' \cap \overline{S'} = \emptyset$. We arrive at the following problem:

Disjoint-FVS	
Input:	graph $G = (V, E)$, integer k, feedback vertex set S of size $k+1$ of G
Output:	a feedback vertex set S^* of G with $|S^*| \leq k$ and $S^* \cap S = \emptyset$, if one exists

If Disjoint-FVS can be solved in $O^*(d^k)$ time, then Comp-FVS can be solved in

$$O^*\left(\sum_{i=0}^{k+1} \binom{k+1}{i} d^i\right) \subseteq O^*((d+1)^k) \text{ time.}$$

Algorithm for Disjoint-FVS
Denote $A := V \setminus S$.



Simplification rules for Disjoint-FVS
Start with $S^* = \emptyset$.

(cycle-in-S)
If $G[S]$ is not acyclic, then return No.

(budget-exceeded)
If $k < 0$, then return No.

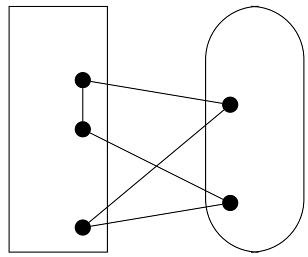
(finished)
If $G - S^*$ is acyclic, then return S^*.

(creates-cycle)
If $\exists v \in A$ such that $G[S \cup \{v\}]$ is not acyclic, then add v to S^* and remove v from G.

(Degree-(≤ 1))
If $\exists v \in V$ with $d_G(v) \leq 1$, then remove v from G.

(Degree-2)
If $\exists v \in V$ with $d_G(v) = 2$ and at least one neighbor of v is in A, then add an edge between the neighbors of v (even if there was already an edge) and remove v from G.

Simplified instance:



Branching rule for Disjoint-FVS
Select a vertex $v \in A$ with at least 2 neighbors in S. Such a vertex exists if no simplification rule applies (for example, we can take a leaf in $G[A]$). Branch into two subproblems:

$v \in S^*$**:** add v to S^*, remove v from G, and decrease k by 1

$v \notin S^*$**:** add v to S

Exercise: Running time

- Prove that this algorithm has running time $O^*(4^k)$.

Hint: Use the measure $k + cc(S)$, where $cc(S)$ is the number of connected components of $G[S]$.

Result for Feedback Vertex Set

Theorem 1. FEEDBACK VERTEX SET *can be solved in* $O^*(5^k)$ *time.*

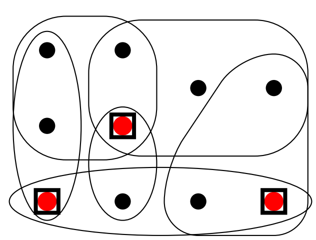
3 Min r-Hitting Set

A *set system* $\mathcal{S}$ is a pair (V, H), where V is a finite set of elements and H is a set of subsets of V. The *rank* of $\mathcal{S}$ is the maximum size of a set in H, i.e., $\max_{Y \in H} |Y|$.

A *hitting set* of a set system $\mathcal{S} = (V, H)$ is a subset X of V such that X contains at least one element of each set in H, i.e., $X \cap Y \neq \emptyset$ for each $Y \in H$.

(universe)-MIN-r-HITTING SET (r-HS)

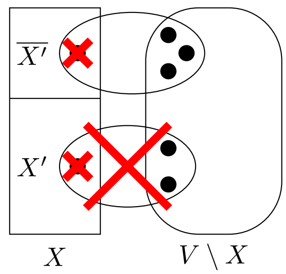
Input:	A rank r set system $\mathcal{S} = (V, H)$
Parameter:	$n = \lvert V \rvert$
Output:	A smallest hitting set of $\mathcal{S}$



Note: The corresponsing decision problem is trivially FPT.

Compression Step

COMP-r-HS Input: set system $\mathcal{S} = (V, H)$, integer k, hitting set X of size $k+1$ of $\mathcal{S}$ Output: a hitting set X^* of size $\leq k$ of $\mathcal{S}$ if one exists



Go over all partitions $(X', \overline{X'})$ of X such that $|X'| \geq 2|X| - n - 1$. Reject a partition if there is a $Y \in H$ such that $Y \subseteq \overline{X'}$. Compute a hitting set X'' of size $\leq k - |X'|$ for (V', H'), where $V' = V \setminus X$ and $H' = \{Y \cap V : Y \in H \wedge Y \cap X' = \emptyset\}$, if one exists. If one exists, then return $X^* = X' \cup X''$.

- The algorithm considers only partitions into $(X', \overline{X'})$ such that $|X'| \geq 2|X| - n - 1$. Number of partitions:

$$O\left(\max\left\{2^{2n/3}, \max_{2n/3 \leq j \leq n} \binom{j}{2j-n}\right\}\right) = O\left(\max_{2n/3 \leq j \leq n} \binom{j}{2j-n}\right)$$

- The subinstances (V', H') where $V' = V \setminus X$ and $H' = \{Y \cap V : Y \in H \wedge Y \cap X' = \emptyset\}$ are instances of $(r-1)$-HS
- Suppose $(r-1)$-HS can be solved in $O^*((\alpha_{r-1})^n)$ time. Then, r-HS can be solved in

$$O^*\left(\max_{2n/3 \leq j \leq n} \binom{j}{2j-n} (\alpha_{r-1})^{n-j}\right) \text{ time} \tag{1}$$

- For example, using a $O(1.6278^n)$ algorithm for 3-HS [Wahlström '07], we obtain a $O(1.8704^n)$ time algorithm for 4-HS [1]3-¿the maximum in (1) is obtained for $j \approx 0.6824 \cdot n$.

Iteration Step

- (V, H) instance of r-HS with $V = \{v_1, v_2, \ldots, v_n\}$
- $V_i = \{v_1, v_2, \ldots, v_i\}$ for $i = 1$ to n
- $H_i = \{Y \in H : Y \subseteq V_i\}$
- Note that $|X_{i-1}| \leq |X_i| \leq |X_{i-1}| + 1$ where X_j is a minimum hitting set of the instance (V_i, H_i)

Theorem 2 ([Fomin, Gaspers, Kratsch, Liedloff, and Saurabh, 2010])**.** *4-HS can be solved in $O(1.8704^n)$ time.*

- One can generalize this result to the counting version of r-HS for any fixed r: count the number of minimum hitting sets of the given set system.

[1] ¡

#r-Hitting Set

Theorem 3 ([Fomin, Gaspers, Kratsch, Liedloff, and Saurabh, 2010])**.** *If there exists a* $O^*((\alpha_{k-1})^n)$ *time algorithm for* $\#(r-1)$*-HS with* $\alpha_{r-1} \leq 2$*, then* $\#r$*-HS can be solved in time*

$$O^*\left(\max_{2n/3 \leq j \leq n}\left\{\binom{j}{2j-n}(\alpha_{r-1})^{n-j}\right\}\right).$$

- If $\alpha_{r-1} \geq 1.6553$, then the following result is better

Theorem 4 ([Fomin, Gaspers, Kratsch, Liedloff, and Saurabh, 2010])**.** *If there exists a* $O^*((\alpha_{r-1})^n)$ *time algorithm for* $\#(r-1)$-HS *with* $\alpha_{k-1} \leq 2$*, then* $\#r$-HS *can be solved in time*

$$\min_{0.5 \leq \beta \leq 1} \max\left\{O^*\left(\binom{n}{\beta n}\right),\ O^*\left(2^{\beta n}(\alpha_{r-1})^{n-\beta n}\right)\right\}.$$

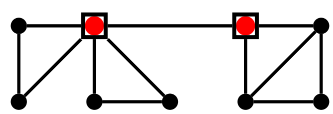
Results for r-HS and #r-HS

r	$\#r$-HS	r-HS
2	$O(1.2377^n)$ [Wahlström '08]	$O(1.2002^n)$ [Xiao, Nagamoshi '13]
3	$O(1.7198^n)$ [Theorem 3]	$O(1.6278^n)$ [Wahlström '07]
4	$O(1.8997^n)$ [Theorem 4]	$O(1.8704^n)$ [Theorem 3]
5	$O(1.9594^n)$ [Theorem 4]	$O(1.9489^n)$ [Theorem 4]
6	$O(1.9824^n)$ [Theorem 4]	$O(1.9781^n)$ [Theorem 4]
7	$O(1.9920^n)$ [Theorem 4]	$O(1.9902^n)$ [Theorem 4]

Exercise

A *cluster graph* is a graph where every connected component is a complete graph.

Cluster Vertex Deletion
Input: Graph $G = (V, E)$, integer k
Parameter: k
Question: Is there a set of vertices $S \subseteq V$ with $|S| \leq k$ such that $G - S$ is a cluster graph?



Recall that G is a cluster graph iff G contains no induced P_3.

- Design an $O^*(2^k)$ time algorithm for Cluster Vertex Deletion.

Hints: (1) Show that the disjoint version of the problem can be solved in polynomial time: given $(G = (V, E), S, k)$ such that $|S| = k + 1$ and $G - S$ is a cluster graph, find a $S^* \subseteq V \setminus S$ with $|S^*| \leq k$ such that $G - S^*$ is a cluster graph. (2) Simplification rule for $v \in V \setminus S$ inducing a P_3 with 2 vertices in S. Reduce to maximum weight matching.

Solution sketch

Disjoint-CVD
Input: graph $G = (V, E)$, integer k, cluster vertex deletion set S of size $k + 1$ of G
Output: a cluster vertex deletion set S^* of G with $|S^*| \leq k$ and $S^* \cap S = \emptyset$, if one exists

Simplification rules:

- If $G[S]$ contains an induced P_3, then return No.
- If $\exists v \in V \setminus S$ such that $G[S \cup \{v\}]$ contains an induced P_3, then set $G \leftarrow G - v$ and $k \leftarrow k - 1$.

Now each vertex in $V \setminus S$ has either no neighbor in S or is adjacent to all the vertices of exactly one cluster of $G[S]$. Reduce the problem to maximum weighted matching in an auxiliary graph where one independent set corresponds to the clusters in $G[S]$ and each vertex in the other independent set corresponds to cliques neighboring exactly one cluster in $G[S]$. It remains to define the edges of the auxiliary graph and their weights.

4 Further Reading

- Chapter 4, *Iterative Compression* in Marek Cygan, Fedor V. Fomin, Łukasz Kowalik, Daniel Lokshtanov, Dániel Marx, Marcin Pilipczuk, MichałPilipczuk, and Saket Saurabh. Parameterized Algorithms. Springer, 2015.
- Section 11.3, *Iterative Compression* in Rolf Niedermeier. Invitation to Fixed Parameter Algorithms. Oxford University Press, 2006.
- Section 6.1, *Iterative Compression: The Basic Technique* in Rodney G. Downey and Michael R. Fellows. Fundamentals of Parameterized Complexity. Springer, 2013.
- Section 6.2, *Edge Bipartization* in Rodney G. Downey and Michael R. Fellows. Fundamentals of Parameterized Complexity. Springer, 2013.