

Glossary

COMP6741: Parameterized and Exact Computation

2015, Semester 2

Glossary

acyclic: A graph is acyclic if it has no **cycle** as a **subgraph**.

bipartite: A graph $G = (V, E)$ is bipartite if its vertex set can be partitioned into two **independent sets**. A partition (A, B) of V into independent sets is called a bipartition of G . The graph G is then often denoted by $G = (A \uplus B, E)$.

Boolean formula: A Boolean formula is constructed from Boolean variables that can take the values true and false (or 1 and 0) by the following operations: conjunction (AND, \wedge), disjunction (OR, \vee), and negation (NOT, \neg).

clique: A subset of vertices $S \subseteq V$ of a graph $G = (V, E)$ is a clique in G if $G[S]$ is a **complete** graph.

closed neighborhood: The closed neighborhood of a vertex v in a graph G is $N_G[v] := \{v\} \cup N_G(v)$. The subscript may be omitted if G is clear from the context.

closed set neighborhood: The closed neighborhood of a subset of vertices $S \subseteq V$ in a graph $G = (V, E)$ is $N_G[S] := \bigcup_{v \in S} N_G[v]$. The subscript may be omitted if G is clear from the context.

coloring: A coloring of a graph $G = (V, E)$ is a function from V to a set of colors (integers) such that every two adjacent vertices in G are mapped to different colors. A k -coloring is a coloring using exactly k colors.

complete: A graph G is complete if there is an edge between each pair of vertices in G . A complete graph on n vertices is denoted by K_n .

Conjunctive Normal Form: A **Boolean formula** is in Conjunctive Normal Form if it is a conjunction of clauses, each clause is a disjunctions of literals, and each literal is a Boolean variable or its negation..

connected: A graph G is connected if there is a **walk** between every two vertices of G .

connected component: **Maximal connected subgraph**.

cycle: **2-regular connected** graph. A cycle on n vertices is denoted C_n .

degree: The degree of a vertex v in a graph G is $d_G(v) := |N_G(v)|$. The subscript may be omitted if G is clear from the context. The degree of a vertex v in a **multigraph:** G is the number of times v appears as an end point of an edge in E .

directed acyclic graph: A directed acyclic graph (DAG) is a **directed graph** that contains no **directed cycle** as a directed subgraph.

directed cycle: **Orientation** of a **cycle** where each vertex has in-degree 1.

directed graph: A directed graph G is an ordered pair (V, A) of a set V of vertices and a set A of arcs, where A is a set of ordered pairs of vertices. Its vertex set is $V(G) = V$ and its arc set is $A(G) = A$.

directed path: **Orientation** of a **path** where each vertex has **in-degree** 1, except the start vertex, which has **in-degree** 0 and **out-degree** 1.

distance: In a graph G , the distance between two vertices $u \in V$ and $v \in V$ is the length of the shortest **walk** minus one between u and v , that is the minimum number of edges needed to be traversed to reach v from u and it is denoted by $dist_G(u, v)$.

dominating set: A subset of vertices $S \subseteq V$ of a graph $G = (V, E)$ is a dominating set of G if $N_G[S] = V$.

feedback vertex set: A feedback vertex set of a graph $G = (V, E)$ is a subset of vertices $S \subseteq V$ such that $G - S$ is acyclic.

forest: **Acyclic** graph.

graph: A (simple, undirected) graph G is an ordered pair (V, E) of a set V of vertices and a set E of edges, where E is a set of unordered pairs of distinct vertices. Its vertex set is $V(G) = V$ and its edge set is $E(G) = E$.

in-degree: The in-degree of a vertex v in a directed graph D is $d_D^-(v) := |\{uv \in A\}|$. The subscript may be omitted if D is clear from the context.

independent set: A subset of vertices $S \subseteq V$ of a graph $G = (V, E)$ is an independent set of G if $G[S]$ has no edges.

induced subgraph: For a graph $G = (V, E)$ and a vertex set $S \subseteq V$, the subgraph of G induced on S is the graph $G[S] := (S, \{uv \in E : u, v \in S\})$.

maximal (set): For a set \mathcal{S} of subsets of a ground set U , a set $X \in \mathcal{S}$ is maximal if there exists no set $Y \in \mathcal{S}$ with $X \subsetneq Y$.

maximum (set): For a set \mathcal{S} of subsets of a ground set U , a set $X \in \mathcal{S}$ is maximum if there exists no set $Y \in \mathcal{S}$ with $|Y| > |X|$.

maximum degree: The maximum degree of a graph $G = (V, E)$ is $\Delta(G) := \max_{v \in V} d_G(v)$.

minimum degree: The minimum degree of a graph $G = (V, E)$ is $\delta(G) := \min_{v \in V} d_G(v)$.

multigraph: A multigraph G is an ordered pair (V, E) of a set V of vertices and a *multiset* E of edges, where E is a multiset of unordered pairs of vertices. Its vertex set is $V(G) = V$ and its edge set is $E(G) = E$.

open neighborhood: The (open) neighborhood of a vertex v in a graph $G = (V, E)$ is $N_G(v) := \{u \in V : uv \in E\}$. The subscript may be omitted if G is clear from the context.

open set neighborhood: The (open) neighborhood of a subset of vertices $S \subseteq V$ in a graph $G = (V, E)$ is $N_G(S) := N_G[S] \setminus S$. The subscript may be omitted if G is clear from the context.

order of growth: Let $g : \mathbb{R}_+ \rightarrow \mathbb{R}_+$ be a function. The set $O(g(n))$ contains every function f such that there exist $c, n_0 \geq 0$ such that $f(n) \leq c \cdot g(n)$ for every $n \geq n_0$. The set $o(g(n))$ contains every function f such that for every $\epsilon > 0$ there exists a $n_0 \geq 0$ such that $f(n) \leq \epsilon \cdot g(n)$ for every $n \geq n_0$. For the set $\Omega(g(n))$, we have that $f(n) \in \Omega(g(n))$ iff $g(n) \in O(f(n))$. For the set $\omega(g(n))$, we have that $f(n) \in \omega(g(n))$ iff $g(n) \in o(f(n))$. For the set $\Theta(g(n))$, we have that $\Theta(g(n)) = O(g(n)) \cap \Omega(g(n))$.

orientation: An orientation of a graph G is a directed graph D that has exactly one arc for each edge of G with the same endpoints.

out-degree: The out-degree of a vertex v in a directed graph D is $d_D^+(v) := |\{vu \in A\}|$. The subscript may be omitted if D is clear from the context.

path: Tree with maximum degree 2. A path on n vertices is denoted P_n .

regular: A graph is d -regular if each of its vertices has degree d . A graph is regular if it is d -regular for some d .

subgraph: A graph H is a subgraph of a graph G if $V(H) \subseteq V(G)$ and $E(H) \subseteq E(G)$.

tree: Acyclic, connected graph.

vertex cover: A subset of vertices $S \subseteq V$ of a graph $G = (V, E)$ is a vertex cover of G if each edge of G is incident to at least one vertex of S .

vertex removal: For a graph $G = (V, E)$ and a vertex set $S \subseteq V$, the graph obtained by removing S from G is $G - S := G[V \setminus S]$. If $S = \{u\}$, we may write $G - u$ instead of $G - \{u\}$.

walk: Sequence of vertices in a graph, with each vertex being adjacent to the vertices immediately preceding and succeeding it in the sequence.

Problem Definitions

k -COLORING

Given a graph G , determine if there is a coloring of G with at most k colors.

k -SAT

Given a Boolean formula in Conjunctive Normal Form where each clause has at most k literals, determine if there is an assignment of its variables such that the formula evaluates to true.

DOMINATING SET

Given a graph G and an integer k , determine whether G has a dominating set of size k .

FEEDBACK VERTEX SET

Given a (multi)graph G and an integer k , determine whether G has a feedback vertex set of size at most k .

INDEPENDENT SET

Given a graph G and an integer k , determine whether G has an independent set of size k .

MAXIMUM INDEPENDENT SET

Given a graph G , find an independent set of G of maximum cardinality.

MINIMUM VERTEX COVER

Given a graph G , find a **vertex cover** of G of minimum cardinality.

SAT

Given a Boolean formula, determine if there is an assignment of its variables such that the formula evaluates to true.

TRAVELING SALESMAN PROBLEM

Given a set $\{1, \dots, n\}$ of n cities, the distance $d(i, j)$ between every two cities i and j , and an integer k , determine whether there is a tour with total distance at most k . A *tour* is a permutation of the cities starting and ending in city 1.

VERTEX COVER

Given a graph G and an integer k , determine whether G has a **vertex cover** of size k .