

**Welcome!**

**COMP1511 18s1**

**Programming Fundamentals**

# COMP1511 18s1

## — Lecture 21 —

### Analysis, Sorting, Searching

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## Don't panic!

**assignment 3** out now!

week 11's tute/lab help you get started

week 11

**lab** solutions now out **weekly test** due **friday**

don't forget about **help sessions!**

see course website for details

there **will be** lectures next week!

(week 13)

# Questions?

unanswered questions?

ask on **Ed**

[edstem.org/courses/1950/](https://edstem.org/courses/1950/)

...

<https://echo360.org.au/>

note: you may need to go via **Moodle**

<https://moodle.telt.unsw.edu.au>

(let me know if you can/can't access it!)

**What topics are you confused about? What questions do you have?**

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What is your response?

# COMP1511 vs future COMP courses

## COMP1511

gets you thinking like a **programmer**  
solving problems by developing programs  
expressing your ideas in the C language

# COMP1511 vs future COMP courses

## future COMP courses

(e.g. COMP2521)

gets you thinking like a **computer scientist**

knowing fundamental techniques/structures

able to reason about applicability/**effectiveness**

able to **analyse** behaviour/correctness of programs

# COMP1511 vs future COMP courses

for now... just a taster

(you'll have to take COMP2521 for more!)



# introducing: **analysis**

putting the **science** in **computer science**

for when “it works!” isn’t good enough

# what makes software **good**?

# what makes software **good**?

**correctness?**

# what makes software **good**?

correctness?

**efficiency?**

# what makes software **good**?

correctness?

efficiency?

**clear, maintainable code?**

# what makes software **good**?

correctness?

efficiency?

clear, maintainable code?

**usability?**

today: **efficiency**

# Efficiency

COMP1511 focuses on writing **correct** programs

but

**efficiency** is also important

often need to consider:

**execution time**

**memory use**

a **correct** but **too slow** program can be useless



# Efficiency

efficiency often depends on the **size** of the data being processed

understanding this dependency lets us  
**predict** program performance  
on larger data

....

informal exploration in COMP1511 - much more in COMP2521 and COMP3121

# Analysis of Algorithms

how can we find out whether a program is efficient or not?

**empirical approach** - run the program  
several times with different input sizes  
and measure the time taken

**theoretical approach** - try to count the number of  
"operations" performed by the algorithm  
on input of size  $n$

# Linear Search Unordered Array - Code

```
int linear_search(int array[], int length, int x) {  
    for (int i = 0; i < length; i = i + 1) {  
        if (array[i] == x) {  
            return 1;  
        }  
    }  
    return 0;  
}
```

# Linear Search Unordered Array - Informal Analysis

Operations:

- start at first element
- inspect each element in turn
- stop when find **X** or reach end

If there are **N** elements to search:

- best case check 1 element
- worst case check N elements
- if in list on average check  $N/2$  elements
- if not in list check N elements

# Linear Search Ordered Array - Code

```
int linear_ordered(int array[], int length, int x) {  
    for (int i = 0; i < length; i = i + 1) {  
        if (array[i] == x) {  
            return 1;  
        } else if (array[i] > x) {  
            return 0;  
        }  
    }  
    return 0;  
}
```

# Linear Search Ordered Array - Informal Analysis

20

Operations:

- start at first element
- inspect each element in turn
- stop when find **X** or find value **X** or reach end

If there are **N** elements to search:

- best case check 1 element
- worst case check N elements
- if in list on average check  $N/2$  elements
- if not in list on average check  $N/2$  elements

# Binary Search Ordered Array - Code

```
int binary_search(int array[], int length, int x) {
    int lower = 0;
    int upper = length - 1;
    while (lower <= upper) {
        int mid = (lower + upper) / 2;
        if (array[mid] == x) {
            return 1;
        } else if (array[mid] > x) {
            upper = mid - 1;
        } else {
            lower = mid + 1;
        }
    }
    return 0;
}
```

# Binary Search Ordered Array - Informal Analysis

Operations:

- start with entire array
- at each step halve the range the element may be in
- stop when find **X** or range is empty

If there are **N** elements to search

- best case check 1 element
- worst case check  $\log_2(N)+1$  elements
- if in list on average check  $\sim \log_2(N)$  elements



# Binary Search Ordered Array - Informal Analysis

$\log_2(N)$  grows very slowly:

- $\log_2(10) = 3.3$
- $\log_2(1000) = \sim 10$
- $\log_2(1000000) = \sim 20$
- $\log_2(1000000000) = \sim 30$
- $\log_2(1000000000000) = \sim 40$

Physicists estimate  $10^{80}$  atoms in universe:  $\log_2(10^{80}) = 240$ [1ex]

Binary search all atoms in universe in  $< 1$  microsecond

# let's look at: **sorting**

# Sorting

**sort:** rearrange a sequence so it is in non-decreasing order

## why?

sorted sequence can be searched efficiently  
items with equal keys are located together

## why not?

simple obvious algorithms too slow to sort large sequences  
(better algorithms can sort very large sequences)

# Sorting Algorithms

there are **many** different sorting algorithms

we'll look at one **slow** obvious algorithm:

**bubblesort**

and at one **fast** algorithm:

**quicksort**

(SortVis: <https://sorting.alhinds.com>)

# Bubblesort

go through the array, comparing pairs of elements

**swap** the elements if they're in the wrong order

...

repeat until sorted

# Bubblesort

```
// our array of numbers
3  1  4  9  5

// compare the first pair
[3] [1] 4  9  5

// they're in the wrong order, so swap
[1] [3] 4  9  5

// compare the second pair
1  [3] [4] 9  5

// compare the third pair
1  3  [4] [9] 5

// compare the fourth pair
1  3  4  [9] [5]

// they're in the wrong order, so swap
1  3  4  [5] [9]
```

# Bubblesort - Code

```
void bubblesort(int array[], int length) {
    int swapped = 1;
    while (swapped) {
        swapped = 0;
        for (int i = 1; i < length; i = i + 1) {
            if (array[i] < array[i - 1]) {
                int tmp = array[i];
                array[i] = array[i - 1];
                array[i - 1] = tmp;
                swapped = 1;
            }
        }
    }
}
```

# Quicksort

faster than bubblesort

**divide and conquer**

(make the problem smaller each time)

works by dividing the array into two smaller arrays  
then sorting the two smaller arrays

...

it does this by choosing a **pivot**  
then moving all of the **smaller** elements to its left and  
all of the **larger** elements to its right



# Quicksort - Code

```
void quicksort(int array[], int length) {
    quicksort1(array, 0, length - 1);
}

void quicksort1(int array[], int lo, int hi) {
    if (lo >= hi) {
        return;
    }
    int p = partition(array, lo, hi);
    // sort lower part of array
    quicksort1(array, lo, p);
    // sort upper part of array
    quicksort1(array, p + 1, hi);
}
```

# Quicksort Partition - Code

```
int partition(int array[], int lo, int hi) {
    int i = lo, j = hi;
    int pivotValue = array[(lo + hi) / 2];
    while (1) {
        while (array[i] < pivotValue) {
            i = i + 1;
        }
        while (array[j] > pivotValue) {
            j = j - 1;
        }
        if (i >= j) {
            return j;
        }
        int temp = array[i];
        array[i] = array[j];
        array[j] = temp;
        i = i + 1;
        j = j - 1;
    }
    return j;
}
```

# Quicksort and Bubblesort Compared

If we instrument quicksort and bubble sort code, we see:

- bubblesort is proportional to  $n^2$
- quicksort is proportional to  $n \log n$
- if  $n$  is small, little difference
- if  $n$  is large, huge difference
- for large  $n$ , you need a good sorting algorithm like quicksort