## COMP4418: Knowledge Representation and Reasoning-Solutions to Exercise 1 Propositional Logic

1. (i) $(\neg J a \wedge \neg J o) \rightarrow T$

Where:
Ja: Jane is in town
Jo: John is in town
$T$ : we will play tennis
(ii) $R \vee \neg R$

Where:
$R$ : it will rain today
(iii) $\neg S \rightarrow \neg P$

Where:
$S$ : you study
$P$ : you will pass this course
(iv) $D \rightarrow(B \vee S)$

Where:
$D: I$ ate dinner
B: I drink bubble tea
S: I drink soft drink
(v) $(V \wedge D) \rightarrow(R \wedge \neg F)$

Where:
$V: 80 \%$ of adults are fully vaccinated
D: COVID-19 cases begin to drop
$R$ : lockdown restrictions are eased
$F$ : international flights immediately resume
2. (i) $P \rightarrow Q$
$\neg P \vee Q($ remove $\rightarrow)$
(ii) $(P \rightarrow \neg Q) \rightarrow R$
$\neg(\neg P \vee \neg Q) \vee R$ (remove $\rightarrow$ )
$(\neg \neg P \wedge \neg \neg Q) \vee R$ (De Morgan)
$(P \wedge Q) \vee R$ (Double Negation)
$(P \vee R) \wedge(Q \vee R)($ Distribute $\vee$ over $\wedge)$
(iii) $\neg(P \wedge \neg Q) \rightarrow(\neg R \vee \neg Q)$
$\neg \neg(P \wedge \neg Q) \vee(\neg R \vee \neg Q)$ (remove $\rightarrow$ )
$(P \wedge \neg Q) \vee(\neg R \vee \neg Q)$ (Double Negation)
$(P \vee \neg R \vee \neg Q) \wedge(\neg Q \vee \neg R \vee \neg Q)$ (Distribute $\vee$ over $\wedge$ )
This can be further simplified to: $((P \vee \neg R \vee \neg Q) \wedge(\neg Q \vee \neg R)$
And in fact this cab be simplified to $\neg Q \vee \neg R$ since $(\neg Q \vee \neg R) \vdash$ $(P \vee \neg R \vee \neg Q)$
(iv) $(\neg P \rightarrow Q) \rightarrow(Q \rightarrow \neg R)$
$\neg(\neg P \rightarrow Q) \vee(Q \rightarrow \neg R)($ remove $\rightarrow)$

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\(\neg(P \vee Q) \vee(\neg Q \vee \neg R)\) (remove \(\rightarrow\) )
\((\neg P \wedge \neg Q) \vee(\neg Q \vee \neg R)\) (De Morgan)
\((\neg P \vee \neg Q \vee \neg R) \wedge(\neg Q \vee \neg R)(\) Distribute \(\vee\) over \(\wedge)\)
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(v) $\neg(\neg P \vee Q) \vee(\neg R \rightarrow S)$
$\neg(\neg P \vee Q) \vee(\neg \neg R \vee S)$ (remove $\rightarrow$ )
$(\neg \neg P \wedge \neg Q) \vee(\neg \neg R \vee S)$ (De Morgan)
$(P \wedge \neg Q) \vee(R \vee S)$ (Double Negation)
$(P \vee R \vee S) \wedge(\neg Q \vee R \vee S)($ Distribute $\vee$ over $\wedge)$
3. (i)

| $P$ | $Q$ | $P \rightarrow Q$ | $\neg Q$ | $\neg P$ |
| :--- | :--- | :--- | :--- | :--- |
| $T$ | $T$ | $T$ | $F$ | $F$ |
| $T$ | $F$ | $F$ | $T$ | $F$ |
| $F$ | $T$ | $T$ | $F$ | $T$ |
| $F$ | $F$ | $T$ | $T$ | $T$ |

In all rows where both $P \rightarrow Q$ and $\neg Q$ are true, $\neg P$ is also true. Therefore, inference is valid.
(ii)

| $P$ | $Q$ | $\neg P$ | $\neg Q$ | $P \rightarrow Q$ | $\neg Q \rightarrow \neg P$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $T$ | $T$ | $F$ | $F$ | $T$ | $T$ |
| $T$ | $F$ | $F$ | $T$ | $F$ | $F$ |
| $F$ | $T$ | $T$ | $F$ | $T$ | $T$ |
| $F$ | $F$ | $T$ | $T$ | $T$ | $T$ |

In all rows where $P \rightarrow Q$ is true, $\neg Q \rightarrow \neg P$ is also true. Therefore, inference is valid.

| $P$ | $Q$ | $R$ | $P \rightarrow Q$ | $Q \rightarrow R$ | $P \rightarrow R$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $T$ | $T$ | $T$ | $T$ | $T$ | $T$ |
| $T$ | $T$ | $F$ | $T$ | $F$ | $F$ |
| $T$ | $F$ | $T$ | $F$ | $T$ | $T$ |
| $T$ | $F$ | $F$ | $F$ | $T$ | $F$ |
| $F$ | $T$ | $T$ | $T$ | $T$ | $T$ |
| $F$ | $T$ | $F$ | $T$ | $F$ | $T$ |
| $F$ | $F$ | $T$ | $T$ | $T$ | $T$ |
| $F$ | $F$ | $F$ | $T$ | $T$ | $T$ |

In all rows where both $P \rightarrow Q$ and $Q \rightarrow R$ are true, $P \rightarrow R$ is also true. Therefore, inference is valid.

| $P$ | $Q$ | $R$ | $Q \wedge R$ | $P \rightarrow Q$ | $P \rightarrow R$ | $P \rightarrow(Q \wedge R)$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $T$ | $T$ | $T$ | $T$ | $T$ | $T$ | $T$ |
| $T$ | $T$ | $F$ | $F$ | $T$ | $F$ | $F$ |
| $T$ | $F$ | $T$ | $F$ | $F$ | $T$ | $F$ |
| $T$ | $F$ | $F$ | $F$ | $F$ | $F$ | $F$ |
| $F$ | $T$ | $T$ | $T$ | $T$ | $T$ | $T$ |
| $F$ | $T$ | $F$ | $F$ | $T$ | $T$ | $T$ |
| $F$ | $F$ | $T$ | $F$ | $T$ | $T$ | $T$ |
| $F$ | $F$ | $F$ | $F$ | $T$ | $T$ | $T$ |

In all rows where both $P \rightarrow Q$ and $P \rightarrow R$ are true, $P \rightarrow(Q \wedge R)$ is also true. Therefore, inference is valid.
(v)

| $P$ | $Q$ | $R$ | $P \wedge Q$ | $Q \rightarrow R$ | $P \rightarrow(Q \rightarrow R)$ | $(P \wedge Q) \rightarrow R$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $T$ | $T$ | $T$ | $T$ | $T$ | $T$ | $T$ |
| $T$ | $T$ | $F$ | $T$ | $F$ | $F$ | $F$ |
| $T$ | $F$ | $T$ | $F$ | $T$ | $T$ | $T$ |
| $T$ | $F$ | $F$ | $F$ | $T$ | $T$ | $T$ |
| $F$ | $T$ | $T$ | $F$ | $T$ | $T$ | $T$ |
| $F$ | $T$ | $F$ | $F$ | $F$ | $T$ | $T$ |
| $F$ | $F$ | $T$ | $F$ | $T$ | $T$ | $T$ |
| $F$ | $F$ | $F$ | $F$ | $T$ | $T$ | $T$ |

In all rows where $P \rightarrow(Q \rightarrow R)$ is true, $(P \wedge Q) \rightarrow R$ is also true. Therefore, inference is valid.
4. (i) $\operatorname{CNF}(P \rightarrow Q)$
$\equiv \neg P \vee Q$
$\operatorname{CNF}(\neg Q)$
$\equiv \neg Q$
CNF ( $\neg \neg P)$
$\equiv P$ (Double Negation)
Proof:

1. $\neg P \vee Q$ (Hypothesis)
2. $\neg Q \quad$ (Hypothesis)
3. $P$ (Negation of Conclusion)
4. $Q \quad 1,3$ Resloution
5. $\square \quad 2,4$ Resloution
(ii) $\operatorname{CNF}(P \rightarrow Q)$
$\equiv \neg P \vee Q$
$\operatorname{CNF}(\neg(\neg Q \rightarrow \neg P))$
$\equiv \neg(\neg \neg Q \vee \neg P)$ (Remove $\rightarrow$ )
$\equiv \neg(Q \vee \neg P)$ (Double Negation)
$\equiv \neg Q \wedge \neg \neg P$ (De Morgan)
$\equiv \neg Q \wedge P$ (Double Negation)
Proof:

| Proof. |  |  |
| :--- | :--- | :--- |
| 1. | $\neg P \vee Q$ | (Hypothesis) |
| 2. | $\neg Q$ | (Negation of Conclusion) |
| 3. | $P$ | (Negation of Conclusion) |
| 4. | $\neg P$ | 1,2 Resolution |
| 5. | $\square$ | 3,4 Resolution |

(iii) $P \rightarrow Q, Q \rightarrow R \vdash P \rightarrow R$
$\operatorname{CNF}(P \rightarrow Q)$
$\equiv \neg P \vee Q$
$\operatorname{CNF}(Q \rightarrow R)$
$\equiv \neg Q \vee R$
$\operatorname{CNF}(\neg(P \rightarrow R))$
$\equiv \neg(\neg P \vee R)$ (Remove $\rightarrow$ )
$\equiv \neg \neg P \wedge \neg R$ (De Morgan)
$\equiv P \wedge \neg R$ (Double Negation)
Proof:

1. $\neg P \vee Q \quad$ (Hypothesis)
2. $\neg Q \vee R \quad$ (Hypothesis)
3. $P$ (Negation of Conclusion)
4. $\neg R \quad$ (Negation of Conclusion)
5. $Q \quad 1,3$ Resolution
6. $R \quad 2,5$ Resolution
7. 

4, 6 Resolution
(iv) $P \rightarrow Q, P \rightarrow R \vdash P \rightarrow(Q \wedge R)$
$\operatorname{CNF}(P \rightarrow Q)$
$\equiv \neg P \vee Q$
$\operatorname{CNF}(P \rightarrow R)$
$\equiv \neg P \vee R$
$\operatorname{CNF}(\neg(P \rightarrow(Q \wedge R)))$
$\equiv \neg(\neg P \vee(Q \wedge R))$ (Remove $\rightarrow$ )
$\equiv \neg \neg P \wedge \neg(Q \wedge R)$ (De Morgan)
$\equiv P \wedge(\neg Q \vee \neg R)$ (Double Negation, De Morgan)
Proof:

1. $\neg P \vee Q \quad$ (Hypothesis)
2. $\neg P \vee R \quad$ (Hypothesis)
3. $P$ (Negation of Conclusion)
4. $\neg Q \vee \neg R \quad$ (Negation of Conclusion)
5. $Q \quad 1,3$ Resolution
6. $R \quad 2,3$ Resolution
7. $\neg R \quad 4,5$ Resolution
8. 

6, 7 Resolution
(v) $P \rightarrow(Q \rightarrow R) \vdash(P \wedge Q) \rightarrow R$
$\operatorname{CNF}(P \rightarrow(Q \rightarrow R))$
$\equiv P \rightarrow(\neg Q \vee R)$ (Remove $\rightarrow$ )
$\equiv \neg P \vee(\neg Q \vee R)($ Remove $\rightarrow)$
$\equiv \neg P \vee \neg Q \vee R$
$\operatorname{CNF}(\neg((P \wedge Q) \rightarrow R))$
$\equiv \neg(\neg(P \wedge Q) \vee R)$ (Remove $\rightarrow$ )
$\equiv \neg \neg(P \wedge Q) \wedge \neg R$ (De Morgan)
$\equiv(P \wedge Q) \wedge \neg R$ (Double Negation)
$\equiv P \wedge Q \wedge \neg R$

## Proof:

| Proof. |  |  |
| :---: | :--- | :--- |
| 1. | $\neg P \vee \neg Q \vee R$ | (Hypothesis) |
| 2. | $P$ | (Negation of Conclusion) |
| 3. | $Q$ | (Negation of Conclusion) |
| 4. | $\neg R$ | (Negation of Conclusion) |
| 5. | $\neg Q \vee R$ | 1,2 Resolution |
| 6. | $R$ | 3,5 Resolution |
| 7. | $\square$ | 4,6 Resolution |

5. (i) $((P \vee Q) \wedge \neg P) \rightarrow Q$

| $P$ | $Q$ | $\neg P$ | $P \vee Q$ | $(P \vee Q) \wedge \neg P$ | $((P \vee Q) \wedge \neg P) \rightarrow Q$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $T$ | $T$ | $F$ | $T$ | $F$ | $T$ |
| $T$ | $F$ | $F$ | $T$ | $F$ | $T$ |
| $F$ | $T$ | $T$ | $T$ | $T$ | $T$ |
| $F$ | $F$ | $T$ | $F$ | $F$ | $T$ |

Last column is always true no matter what truth assignment to the atoms $P$ and $Q$. Therefore $((P \vee Q) \wedge \neg P) \rightarrow Q$ is a tautology.
(ii) $((P \rightarrow Q) \wedge \neg(P \rightarrow R)) \rightarrow(P \rightarrow Q)$

| $P$ | $Q$ | $R$ | $P \rightarrow Q$ | $\neg(P \rightarrow R)$ | $(P \rightarrow Q) \wedge \neg(P \rightarrow R)$ | $((P \rightarrow Q) \wedge \neg(P \rightarrow R)) \rightarrow(P \rightarrow Q)$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $T$ | $T$ | $T$ | $T$ | $F$ | $F$ | $T$ |
| $T$ | $T$ | $F$ | $T$ | $T$ | $T$ | $T$ |
| $T$ | $F$ | $T$ | $F$ | $F$ | $F$ | $T$ |
| $T$ | $F$ | $F$ | $F$ | $T$ | $F$ | $T$ |
| $F$ | $T$ | $T$ | $T$ | $F$ | $F$ | $T$ |
| $F$ | $T$ | $F$ | $T$ | $F$ | $F$ | $T$ |
| $F$ | $F$ | $T$ | $T$ | $F$ | $F$ | $T$ |
| $F$ | $F$ | $F$ | $T$ | $F$ | $F$ | $T$ |

Last column is always true no matter what truth assignment to the atoms $P, Q$ and $R$. Therefore $((P \rightarrow Q) \wedge \neg(P \rightarrow R)) \rightarrow(P \rightarrow Q)$ is a tautology.
(iii) $\neg(\neg P \wedge P) \wedge P$

| $P$ | $\neg P$ | $\neg P \wedge P$ | $\neg(\neg P \wedge P)$ | $\neg(\neg P \wedge P) \wedge P$ |
| :--- | :--- | :--- | :--- | :--- |
| $T$ | $F$ | $F$ | $T$ | $T$ |
| $F$ | $T$ | $F$ | $T$ | $F$ |

Last column is not always true. Therefore $\neg(\neg P \wedge P) \wedge P$ is not a tautology.
(iv) $(P \vee Q) \rightarrow \neg(\neg P \wedge \neg Q)$

| $P$ | $Q$ | $\neg P$ | $\neg Q$ | $P \vee Q$ | $\neg P \wedge \neg Q$ | $\neg(\neg P \wedge \neg Q)$ | $(P \vee Q) \rightarrow \neg(\neg P \wedge \neg Q)$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $T$ | $T$ | $F$ | $F$ | $T$ | $F$ | $T$ | $T$ |
| $T$ | $F$ | $F$ | $T$ | $T$ | $F$ | $T$ | $T$ |
| $F$ | $T$ | $T$ | $F$ | $T$ | $F$ | $T$ | $T$ |
| $F$ | $F$ | $T$ | $T$ | $F$ | $T$ | $F$ | $T$ |

(v) $(P \vee Q) \wedge \neg(P \wedge Q)$

| $P$ | $Q$ | $P \vee Q$ | $P \wedge Q$ | $\neg(P \wedge Q)$ | $(P \vee Q) \wedge \neg(P \wedge Q)$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $T$ | $T$ | $T$ | $T$ | $F$ | $F$ |
| $T$ | $F$ | $T$ | $F$ | $T$ | $T$ |
| $F$ | $T$ | $T$ | $F$ | $T$ | $T$ |
| $F$ | $F$ | $F$ | $F$ | $T$ | $F$ |

Last column is not always true. Therefore $(P \vee Q) \wedge \neg(P \wedge Q)$ is not a tautology.
6. (i) $\operatorname{CNF}(\neg(((P \vee Q) \wedge \neg P) \rightarrow Q)) \equiv \neg(\neg((P \vee Q) \wedge \neg P) \vee Q)$ (Remove $\rightarrow$ )
$\equiv \neg \neg((P \vee Q) \wedge \neg P) \wedge \neg Q)($ DeMorgan $)$
$\equiv(P \vee Q) \wedge \neg P) \wedge \neg Q$ (Double Negation)

Proof:

1. $\quad P \vee Q \quad$ (Negated Conclusion)
2. $\neg P \quad$ (Negated Conclusion)
3. $\neg Q \quad$ (Negated Conclusion)
4. $Q \quad 1,2$ Resolution
5. 

3, 4 Resolution
Therefore $\neg(((P \vee Q) \wedge \neg P) \rightarrow Q)$ is a tautology.
(ii) $\operatorname{CNF}(\neg(((P \rightarrow Q) \wedge \neg(P \rightarrow R)) \rightarrow(P \rightarrow Q)))$
$\equiv \neg(\neg((\neg P \vee Q) \wedge \neg(\neg P \vee R)) \vee(\neg P \vee Q))$ (Remove $\rightarrow)$
$\equiv \neg \neg((\neg P \vee Q) \wedge \neg(\neg P \vee R)) \wedge \neg(\neg P \vee Q)$ (De Morgan)
$\equiv(\neg P \vee Q) \wedge(\neg \neg P \wedge \neg R) \wedge(\neg \neg P \wedge \neg Q)$ (Double Negation and De Morgan)
$\equiv(\neg P \vee Q) \wedge(P \wedge \neg R) \wedge(P \wedge \neg Q)$ (Double Negation)
Proof:

1. $\neg P \vee Q \quad$ (Negated Conclusion)
2. $P \quad$ (Negated Conclusion)
3. $\neg R \quad$ (Negated Conclusion)
4. $\neg Q \quad$ (Negated Conclusion)
5. $Q \quad$ 1, 2 Resolution
6. 

4,5 Resolution
Therefore $((P \rightarrow Q) \wedge \neg(P \rightarrow R)) \rightarrow(P \rightarrow Q)$ is a tautology.
(iii) $\operatorname{CNF}(\neg(\neg(\neg P \wedge P) \wedge P))$
$\equiv \neg \neg(\neg P \wedge P) \vee \neg P$ (De Morgan)
$\equiv(\neg P \wedge P) \vee \neg P$ (Double Negation)
$\equiv(\neg P \vee \neg P) \wedge(P \vee \neg P)$ (Distribute $\wedge$ over $\vee$ )
$\equiv \neg P$ (Can simplify to this by removing repetition and tautologies)
Proof:

1. $\neg P \quad$ (Negated Conclusion)

Cannot obtain empty clause using resolution so $\neg(\neg P \wedge P) \wedge P$ is not a tautology.
(iv) $\operatorname{CNF}(\neg((P \vee Q) \rightarrow \neg(\neg P \wedge \neg Q)))$
$\equiv \neg(\neg(P \vee Q) \vee \neg(\neg P \wedge \neg Q))$ (Remove $\rightarrow)$
$\equiv \neg \neg(P \vee Q) \wedge \neg \neg(\neg P \wedge \neg Q))$ (De Morgan)
$\equiv(P \vee Q) \wedge(\neg P \wedge \neg Q))$ (Double Negation)

Proof:

1. $(P \vee Q) \quad$ (Negated Conclusion)
2. $\neg Q \quad$ (Negated Conclusion)
3. $\neg P \quad$ (Negated Conclusion)
4. $Q \quad 1,2$ Resolution
5. $\square \quad 3,4$, Resolution

Therefore $(P \vee Q) \rightarrow \neg(\neg P \wedge \neg Q)$ is a tautology.
(v) $\operatorname{CNF}(\neg((P \vee Q) \wedge \neg(P \wedge Q)))$
$\equiv \neg(P \vee Q) \vee \neg \neg(P \wedge Q)$ (De Morgan)
$\equiv \neg(P \vee Q) \vee(P \wedge Q)$ (Double Negation)
$\equiv(\neg P \wedge \neg Q) \vee(P \wedge Q)$ (De Morgan)
$\equiv(\neg P \vee P) \wedge(\neg P \vee Q) \wedge(\neg Q \vee P) \wedge(\neg Q \vee Q)$ (Distribution)
$\equiv(\neg P \vee Q) \wedge(P \vee \neg Q)$ (Removal of tautologies)

Proof:

1. $(\neg P \vee Q) \quad$ (Negated Conclusion)
2. $(P \vee \neg Q) \quad$ (Negated Conclusion)

Cannot obtain empty clause using resolution so $(P \vee Q) \wedge \neg(P \wedge Q)$ is not a tautology.
7. P: I will listen to the album "SOUR" by Olivia Rodrigo

Q: I will watch another episode of The Queen's Gambit
Truth table:
$P \vee Q, \neg Q \models \neg P$

| $P$ | $Q$ | $P \vee Q$ | $\neg Q$ | $\neg P$ |
| :--- | :--- | :--- | :--- | :--- |
| $T$ | $T$ | $T$ | $F$ | $F$ |
| $T$ | $F$ | $T$ | $T$ | $F$ |
| $F$ | $T$ | $T$ | $F$ | $T$ |
| $F$ | $F$ | $F$ | $T$ | $T$ |

This inference is not valid as $\neg P$ is not always true when $(P \vee Q)$ and $\neg Q$ are both true.

Resolution:
$P \vee Q, \neg Q \vdash \neg P$
CNF ( $\neg \neg P)$
$\equiv P$ (Double Negation)

Proof:

1. $\quad P \vee Q \quad$ (Hypothesis)
2. $\neg Q \quad$ (Hypothesis)
3. $P$ (Negated Conclusion)
4. $P \quad 1,2$ Resolution

This inference is not valid as we cannot derive the empty clause using resolution.
8. B: I will drink too much bubble tea

S: I feel sick
Truth table:
$(B \rightarrow S) \vee(S \rightarrow B)$

| $B$ | $S$ | $B \rightarrow S$ | $S \rightarrow B$ | $(B \rightarrow S) \vee(S \rightarrow B)$ |
| :--- | :--- | :--- | :--- | :--- |
| $T$ | $T$ | $T$ | $T$ | $T$ |
| $T$ | $F$ | $F$ | $T$ | $T$ |
| $F$ | $T$ | $T$ | $F$ | $T$ |
| $F$ | $F$ | $T$ | $T$ | $T$ |

Last column is always true. Therefore the statement is a tautology.
Resolution:
$\vdash(B \rightarrow S) \vee(S \rightarrow B)$
$\operatorname{CNF}(\neg((B \rightarrow S) \vee(S \rightarrow B)))$
$\equiv \neg((\neg B \vee S) \vee(\neg S \vee B))$ (Remove $\rightarrow$ )
$\equiv \neg(\neg B \vee S) \wedge \neg(\neg S \vee B)$ (De Morgan)
$\equiv(\neg \neg B \wedge \neg S) \wedge(\neg \neg S \wedge \neg B)$ (De Morgan)
$\equiv(B \wedge \neg S) \wedge(S \wedge \neg B)$ (Double Negation)

Proof:

| 1. | $B \wedge \neg S$ | (Negated Conclusion) |
| :--- | :--- | :--- |
| 2. | $S \wedge \neg B$ | (Negated Conclusion) |
| 3. | $\square$ | (1, 2 Resolution) |

Therefore the sentence is a tautology.

